Sept 17 2021

Errata to Fomin, Lokshtanov, Saurabh and Zehavi "Kernelization" book

The page and statement numbers follows the final (printed) edition by Cambridge.

Many thanks to Frieder Smolny, Satoshi Matsuoka, Julian Schnitzler, Naoyuki Kamiyama,
Tim Jackman, Steve Homer, Jungho Ahn, Danil Sagunov, George Osipov, and Ruben van Niekerk for reporting some of those!

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- page xii, 8th line mulivariate -> multivariate

- page 5, 12th line. Reduction Rule Reduction Rule -> Reduction Rule

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- page 6, In the chess example, the knight moves to e7 instead of f7

- page 15p, 5th line. graphs of girth 5 -> graphs of girth at least 5

- page 16, 19th line. Vertex Cover, Feedback Arc Set in Tournaments, Vertex Cover, Feedback Arc Set in Tournaments (FAST) -> Vertex Cover, Feedback Arc Set in Tournaments (FAST)

- page 16, 20th line. graphs of girth 5 -> graphs of girth at least 5

- page 22 line 4 of Section 2.4 : "The Dominating Set problem it is known..." => "The
Dominating Set problem is known..."

- page 23, 10th line. white vertex dominating more than k black vertices -> white vertex dominating more than $k-|R|$ black vertices

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- page 24: |B| <= k^2 should be |B| <= k(k+1), since one black vertex
covers at most (k+1) black vertices (including itself). The bounds on
W and R+W+B are also affected.

- page 26: In most of reduction rules here, we can use "at least |S|"
instead of "more than |S|"

- page 27, 6th line of the proof of Lemma 2.11) a vertex cover of size $\ell-1$. -> a vertex cover of size at most $\ell-1$.

- page 27, 11th line of the proof of Lemma 2.11) should belong to C. -> should belong to C'.

- pages 29-30: Problem 2.8: Here r is a constant that may be hidden by big-O.

- page 33, Section 3.1 2nd line. "NonBlocker" => "Max Leaf Subtree"

- page 36 line 4: "u and v become leaves": this is not true, since x or y may
coincide with u or v. Here we have to argue that if x coincides with u, then adding xy destroys at most one leaf.

- page 37, 3rd line of Lemma 3.3. The number of subdivided edges is not at most $|L|$, but $|L|+|B|-1\leq 2|L| -3$. This propagates with applications of Lemma 3.3; for instance (3.3) in 44p, 5th line from the bottom of 44p, and 3rd line from the bottom of 46p.

- page 43: FVS.4 "delete it v" => "delete v"

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- page 43: FVS.5 "none-isolated" => "non-isolated"
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- page 43 line 20 "G'-S contain" => "G'-S contains"

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- page 43. Before reduction FVS.5 we need one more reduction rule. Without it rule FVS.5 is not safe.

Lemma: Let v be in T1, and suppose its neighborhood in F is a double edge clique C. Then there exists an optimal FVS S in G that avoids v.

Proof: Let S be an optimal FVS of G. If $v \notin S$ we are done. If v is in S, observe that there is at most one vertex x of C that is is not in S. Let $S' = (S-\{v\}) \cup C$, we have that $|S'| \leq |S|$. We claim that S’ is an FVS. Indeed v has degree at most one in G-S’ and so there are no cycles in G-S’ containing v. Cycles in G-S’ avoiding v also avoid S, contradicting that S is an FVS.

Corollary: Let v in T1 and c be any vertex in C adjacent to v by a double edge. There exists an optimal FVS of G containing c.
Pf: take one that avoids v, it contains c.

This leads to an extra Reduction Rule FVS.4B: if v is in T1 and its neighborhood is a double Edge Clique C, and v is adjacent to c in C by a double Edge, then remove c and reduce the parameter k by 1.

Assuming this rule is applied exhaustively v is not incident to any double edges and the rule may now be applied

Proof of Claim 3.12. (Opposite direction.)

Let v be the vertex of degree 1 in T whose neighborhood in T is a double edge clique. Let u be its neighbor in T. Let w be the vertex resulting from the contraction. We want to prove that if G/uv has a FVS S' of size at most k then G has a FVS of size at most k.
Let S' be an FVS of G/uv.

Case 1: S' does not contain w. Then S' is a vertex set in G, and we claim that S' is a FVS in G. We have that (G-S') / uv = (G/uv) - S'. So (G-S') / uv is acyclic. Contracting an edge (when keeping double edges) can not turn a cyclic graph into acyclic, thus G-S' is acyclic as well.

Case 2: S' contains w. Let S = (S' - {w})\cup {u}. We claim that S is a FVS of G. Suppose not, and let Q be a cycle in G - S. We have that G-S = (G-S’) \cup v. So G-S is an acyclic graph (namely G-S’) plus one vertex v. So Q must go through v. Because of the reduction rule FVS4.a, there are no double edges connecting v with vertices of C. Thus the neighbors of v in Q are two different vertices.
But (since u is in S), both neighbors of v in Q are both in C, contradicting that Q is an induced cycle (since these two neighbors are connected by a double edge in C and thus form a cycle in G and in G'). This is a contradiction.

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- page 45 line 11: "To follows our strategy" => "To follow our strategy"

- page 45: FVS.6 "delete v" => "delete f"

-page 46 line 18. … which endpoints -> whose endpoints

- page 46: Proof of Claim 3.15. “is almost identical to the proof of Claim 3.15”-> Lemma 2.11?

- page 46, 7th line from the bottom. Dince the number of pairs is bounded by $k(k+1)(2k-1)$, then the bound $|T_2|$, $2k(k+1)(2k-1)$ is correct. This propagates to the last two formulas on page 46.

- page 47: Problem 3.3: This problem is very easy since every graph
without isolated vertices has a dominating set of size at most n/2, so
simple kernel with 2k vertices follows. More interesting questions is a 5/3k kerenel, from the original paper.

- page 57, 4th line of Section 4.4. contains a cycle on ℓ vertices. -> contains a cycle on at least $ \ell$ vertices.

- page 58, 2nd line of Problem 4.1. graph G -> a graph G

- page 63, 9th line. from the bottom) $|B|\geq|A|$. -> $|B|\geq q|A|$.

- page 63, 2nd line of the proof of Lemma 5.2 "Next, suppose that that..." => "Next, suppose that..."

- page 72: in Reduction COC.2: "If |B| >=" => "If w(B) >="

- page 73, 11th line. size at least ℓ. -> size at least $\ell+1$.

- page 73: missing qed in the end of the proof of the claim

- page 76, 7th line. allowsus -> allow us

- page 77, in (v): "to the the subtree" => "to the subtree"

- page 79 3rd line: "by more that" => "by more than"

- page 80 item 3 of Reduction FVS12: "Add edges between v and vertices in S such that..." => "Add edges between v and vertices in S in such a way that..."

- page 81 line 14: "such a feedback vertex set the set..." => "such a feedback
vertex set..."

- page 82: For Problem 5.7 a quadratic kernel is easily obtained with
simpler techniques. It is better to ask here for a linear kernel.

- page 91, 10th line of the proof of Theorem 6.5. Conversely, let S be a minimal vertex cover of size k. -> Conversely, let S be a minimal vertex cover of size at most k.

- page 97, 9th line from the bottom. not exist a (resp. strictly) reducible pair -> not exist a (resp. strictly) f-reducible pair

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- page 100, 11th line from the bottom. f-reducible -> f-reducible

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- page 112, 2nd line. $|N_G(A)|\leq|A|$ -> $|N_{G_{mathcal{P}}(A)|\leq|A|$.

- page 112, 3rd line. Define Y as $\delta(\mathcal{P})\setminus A$, instead of $\delta(\mathcal{P})\setminus\{e|e\in A\}$.

- page 112, 4th line the sets of H -> the subsets of $V(H)$

- page 112, 5th line Consider the following partition of H, -> Consider the following partition of $V(H)$,

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- page 113, 12th line flopping -> flipping

- page 113, 6th line from the bottom H is a hypertree -> H has a hypertree

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- page 115, Reduction IST.1 Otherwise, proceed to Rule IST.3. -> Otherwise, proceed to Rule IST.2.

- page 117, Theorem 7.13 ISTadmits -> IST admits

- page 121, 3rd line a an area -> an area

- page 124, 11th line from the bottom Reduction Rule 8.4 -> Reduction Rule SP.1

- page 128, 9th line from the botton $\leq=$ -> $\leq$

- page 129, 3rd line of the proof of Lemma 8.6) Between -> Because

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- page 129, 5th line of the proof of Lemma 8.6 $|F(X)|\leq j$. -> $|F(X)|<j$.

- page 130, 8th line of the proof of Claim 8.8 $|N(w)\cap I|<ik^{i-p-1}$ -> $|N(w)\cap I|<ik^{j-p-1}$

- page 131, 3rd line $I=\bigcap_{x\in S}N(u_x)$, ->

$I=\bigcap_{u_x\in S}N(u_x)$,

- page 131, 17th line Theorem 8.6 -> Lemma 8.6

- page 136, 2nd line from the bottom in the proof Lemma 9.5 -> in the proof of Lemma 9.5

- page 140, 13th line. "than one unaffected clique" -> "than one unaffected maximal clique."

- page 141, Figure 9.3. both incident to K_1 and K_2 -> both incident to K_1 or K_2.

- page 142, 11th line from the bottom)   Observation 9.3 -> Lemma 9.10

- page   148p, 18th line. "all its children are either leaves or labeld by" -> "each of its children either is a leaf or is labeled by."

- page 150, Lemma 9.22. Let G be graph -> Let G be a graph

- page   154, 5th line of the proof of Lemma 9.28."that do not have both endpoints in M" -> “that neither of its ends are in M."
(*235p, 7th line) What is $\tilde{O}$?

- page 169, 6th line. ... if it forms a spanning forest in ... ->... if it forms a   forest in ...

- page 235, 7th line. $x,k\tilde{O}$ -> $(x,k’)$.

- page 239. The title of section 13.2.1. Should be Planar Connected Vertex Cover (not Planar Cluster Vertex Deletion)

- page 159, 4th line. ket -> let

- page 222, inequality in 10th line, $\geq|T''|+k$ -> $\geq|\mathcal{T}''+k$

- page 234, 8th line. least $2k$ -> least $2k'$

- page 234, 16th line. $\geq2k$ -> $\geq2k'$

- page 237, Theorem 13.1. Then? -> Then

- page 240_7 Let C be a *connected* vertex cover...

- page 241^1 Instead of $|N_2|\leq k$ should be $|N_2|=0$.

- page 260 line 19:   a \in \chi(\partial_T(B)) should be u \in \chi(\partial_T(B))

- page 263, 4th line of Lemma 14.18 and 7th line of the proof of Lemma 14.18: $\chi(b')=\chi(b)$. -> $\chi'(b')=\chi(b)$.

- page 264, 11th line from the bottom. $\chi:V(T)\rightarrow V(G)$ -> $\chi:V(T)\rightarrow 2^{V(G)}$

- page 267, the middle equation. Two $\setminus$ between χ functions are missed.

- page 272, 19th line. $OPT_2=(OPT\setminus V_{b_1})\cup X_2$ -> $OPT_2=(OPT_1\setminus V_{b_2})\cup X_2$.

- page 273, Lemma 14.33. k is the width of the tree decomposition.

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- page 291, Theorem 14.56. Robertson et al. et al. -> Robertson et al.

- page 291, Theorem 14.56. For any G excluding H as a minor -> For any G excluding a planar H as a minor

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- page 292 Theorem 14.58. If G is planar -> If G is a connected planar graph.

- page 307. line 5 of the proof of Lemma 15.13.
$G[\chi(T_{v_i} - M_{i-1})]$ should be
$G[\chi(T_{v_i}) \setminus \chi(M_{i-1})]$
Similar corrections for
      line 7 $G[\chi(T - M_i)]$ -> $G[\chi(T) \setminus \chi(M_i)]$
and
line 15
$G[\chi(T_{v_j} - M_{j-1})]$ should be
$G[\chi(T_{v_j}) \setminus \chi(M_{j-1})]$

- pages 310-312 All \hat symbols over T_a χ like in $(\hat{T_a},\hat{\chi_a})$ should be removed.

- page 313, Theorem 15.22) Dominating Set (vc) on planar graphs -> Dominating Set on planar graphs. "

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- page 315, 4th line. Set admits $2^{o(n)}$ time algorithm -> Set admits no $2^{o(n)}$ time algorithm
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- page   328, 2nd line. linear linear -> linear

- page 371, last line: Should be “All but one of its vertices are in S.”

- page 437, 14th line from the bottom. is the restriction of $\Gamma^-_{t,d}$ -> is the restriction of $\Gamma_{t,d}$.

- page 438, 4th line, it suffices to given -> it suffices to give

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- page 440, line 13: “no more hold” should be “no longer hold”.

- page 440, line 14: “setup” should be “set up”.

- page 441 19th line. instance (I,k) In this chapter -> instance (I,k). In this chapter

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- page 454, Definition 23.17. (fEfficient PSAKS) -> (Efficient PSAKS)

- page 481, Question of Multicolored Biclique. $1\leq i\leq\ell$ -> $1\leq i\leq k$